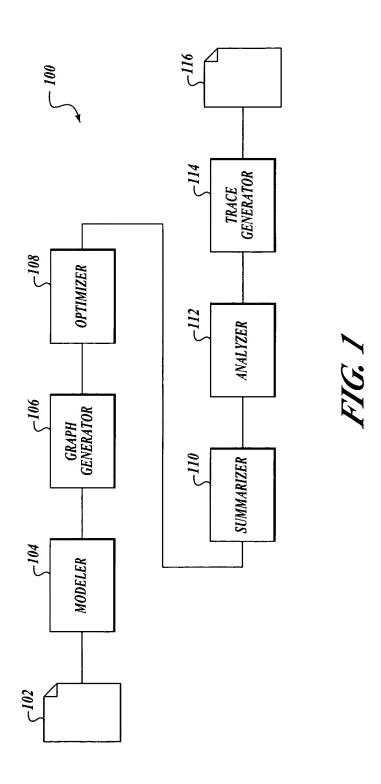


Inventor: Thomas J. Ball et al.
Docket No.: 50037.58US01
Title: METHODS FOR ENHANCING PROGRAM ANALYSIS
Serial No.: 09/866,090



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	[2] if (nU0) { [3] if (?) { [4]; [5] nU0:=F; [6] cE:=T;	} else [7] cE:=T;	[8] if (cE) [9] if (nU0) [10]; else	[11];	B_3 $\left\langle \begin{array}{c} 208 \end{array} \right\rangle$
	[2] if (nU0) { [3] if (?) { [4]; [5] nU0:=F; [6];	} else [7];	[8] if (?) { [9] if (nU0) [10];	[11];	B_2 $\left\langle \begin{array}{c} 206 \end{array} \right\rangle$
void getUnit() {	[2] if (?) { [3] if (?) { [4]; [5]; [6];	} else [7]	[8] if (?) [9] if (?) [10]; else	[11];	B_1 \downarrow 204
int numUnits; int level; void getUnit() { bool canEnter := F;	<pre>if (numUnits = 0) { if (level > 10){ NewUnit(); numUnits := 1; canEnter := T; }</pre>	} else canEnter ;= T;	<pre>if (canEnter) if (numUnits = 0) { assert(F); else</pre>] gotUnit(); }	202
_	<u> </u>		[8] [10]	[11]	Q

FIG. 2

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```
300
                                                                  \( \) 304
        decl g;
                         bebop v1.0: (c) Microsoft Corporation.
                         Done creating bdd variables
        main()
                         Done building transition relations
        begin
                         Label R reachable by following path:
          decl h;
[6]
          h := !g;
[7]
                         Line 12
          A(g,h);
                                                 State g=1 h=0
                         Line 11
[8]
          skip;
                                                 State g=1 h=0
[9]
                         Line 10
          A(g,h);
                                                 State g=1 h=0
[10]
                                Line 22
                                                 State g=1 a1=1 a2=0
          skip;
                                      Line 24
                                                 State g=1 a1=0 a2=1
[11]
          if (g) then
                                      Line 20
        R; skip;
                                                 State g=1 a1=0 a2=1
[12]
                                Line 21
          else
                                                 State g=1 a1=1 a2=0
                                Line 20
[14]
                                                 State g=1 a1=1 a2=0
            skip;
                         Line 9
          fi
                                                 State g=1 h=0
                         Line 8
        end
                                                 State g=1 h=0
                                Line 22
                                                 State g=1 a1=1 a2=0
                                      Line 24
        A(a1,a2)
                                                 State g=1 a1=0 a2=1
                                      Line 20
        begin
                                                 State g=1 a1=0 a2=1
                                Line 21
[20]
          if (a1) then
                                                 State g=1 a1=1 a2=0
                                Line 20
[21]
                                                 State g=1 a1=1 a2=0
            A(a2,a1);
                         Line 7
[22]
            skip;
                                                 State g=1 h=0
                         Line 6
          else
                                                 State g=1
[24]
            g := a2;
          fi
        end
```

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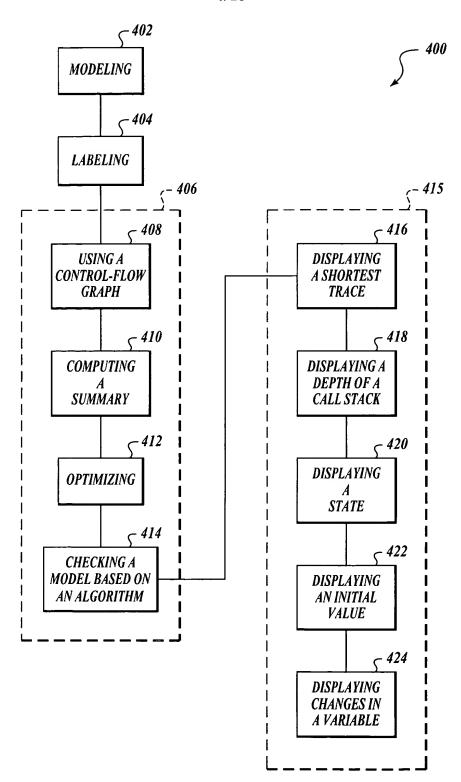


FIG. 4

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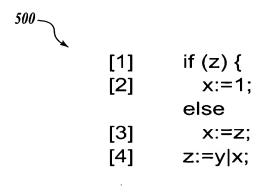


FIG. 5A

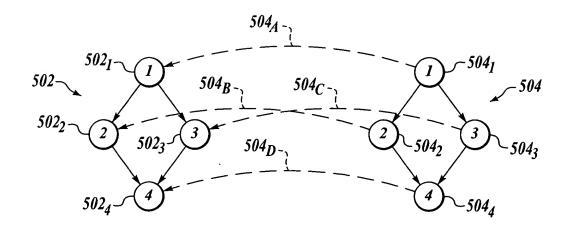
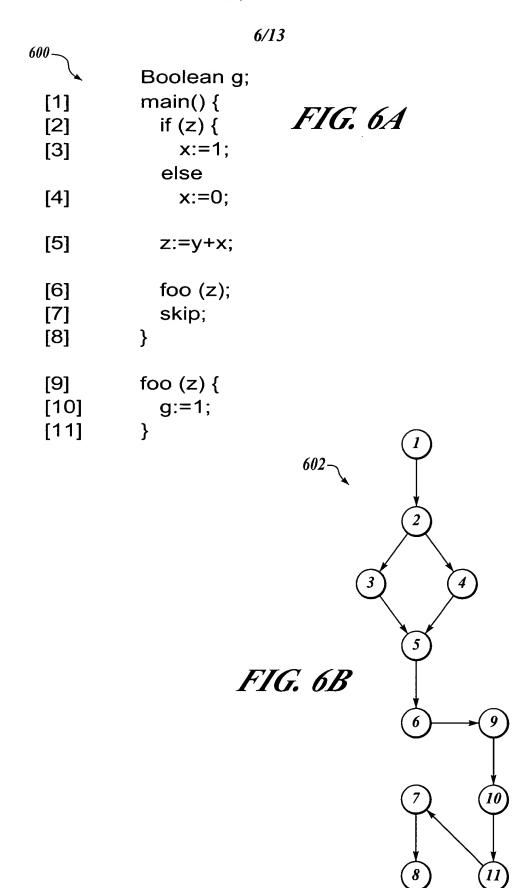


FIG. 5B

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ℐ		
	Λ	$Transfer_{\nu}$
702	skip print goto return	$\lambda\langle\Omega_1,\Omega_2\rangle\cdot(\Omega_2=\Omega_1)$
704	$x_1,, x_k := e_1,, e_k$	$\lambda \langle \Omega_1, \Omega_2 \rangle . (\Omega_2 = \Omega_1 [\mathbf{x}_1 / \Omega_1(\mathbf{e}_1)] [\mathbf{x}_k / \Omega_1(\mathbf{e}_k)])$
706	if(d) while(d) assert(d)	Transfer _{v,true} = $\lambda \langle \Omega_1, \Omega_2 \rangle . ((\Omega_1(d) = 1)^{\wedge} (\Omega_2 = \Omega_1))$ Transfer _{v,false} = $\lambda \langle \Omega_1, \Omega_2 \rangle . ((\Omega_1(d) = 0)^{\wedge} (\Omega_2 = \Omega_1))$
708	$\mathbf{pr}(e_1, \dots, e_k)$	$\lambda \langle \Omega_1, \Omega_2 \rangle . (\Omega_2 = \Omega_1[x_1/\Omega_1(e_1)][x_k/\Omega_1(e_k)]),$ where $x_1,, x_k$ are the formal parameters of pr

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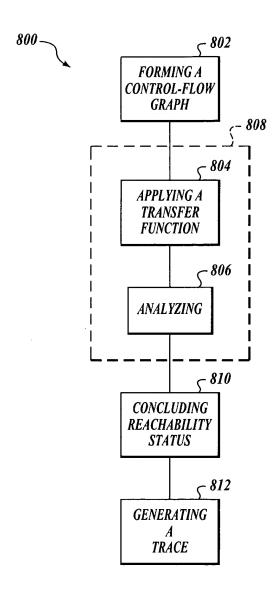


FIG. 8

Title: METHODS FOR ENHANCING PROGRAM ANALYSIS

```
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902 — global
PathEdges, Summary Edges, WorkList
                                                                                           - 900
904 \sim procedure Propagate(v,p)
        begin
   906 \sim if p \neq PathEdges(v) then
     908 \sim PathEdges(v) := PathEdges(v) \cup p
     910 _ Insert v into WorkList fi
912 ← end
914 \sim procedure Reachable(G_B)
  begin for all v \in V_B do PathEdges(v) := \{\}
  917 for all v \in Call_R do SummaryEdges(v) := \{\}
  918 \sim PathEdges(First_R(main)) :=
              \{\langle \Omega, \Omega \rangle \mid \Omega \text{ is any valuation to globals and local variables of main} \}
  920 \sim WorkList := \{First_B(\mathbf{main})\}\
  922 while WorkList \neq \bar{0} do
     924 \sim remove vertex \nu from WorkList
     926 switch (v)
       928 \sim case v \in Call_R
                   Propagate(_{S}^{"}Succ_{B}(v), SelfLoop(Join(PathEdges(v), Transfer_{v})))
                   Propagate(ReturnPt_B(v), Join(PathEdges(v), SummaryEdges(v)))
       934 \sim case v \in Exit_R:
                   for each w \in Succ_R(v) do -936
                      let
                        c \in Call_B such that w=ReturnPt_B(c) and solution 938
                        s = Lift_{\mathbb{C}}(PathEdges(v), ProcOf_{\mathbb{B}}(v)) \sim 940
                     in
                        if s \not\subset SummaryEdges(c) then \checkmark 944
                           SummaryEdges(c) := SummaryEdges(c) \cup s \bigcirc 948
Propagate(w,Join(PathEdges(c),SummaryEdges(c)));
                     ni
       950 \sim case v \in Cond_R:
                   Propagate(Tsucc<sub>B</sub>(v), Join (PathEdges(v), Transfer<sub>v</sub>, true)) - 952
                   Propagate(Fsucc<sub>B</sub>(v), Join (PathEdges(v), Transfer<sub>v</sub>, false)) \sim 954
       956 \sim case v \in V_B - Call_B - Exit_B - Cond_B:
let p = Join(PathEdges(v), Transfer_V) in \sim 958
                     for each w \in Succ_R(v) do \sim 960
                        Propagate(w,p) \sim 962
                   ni
                                                                                 FIG. 9
        end
```

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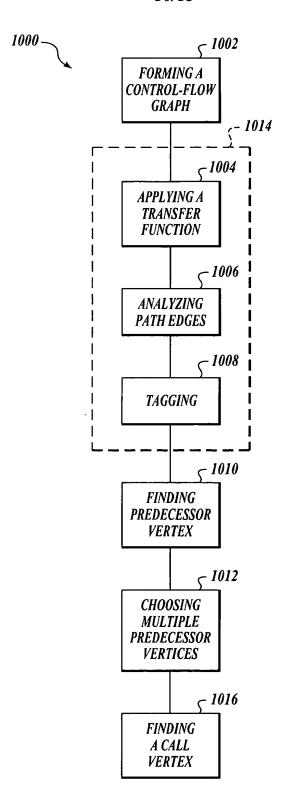


FIG. 10

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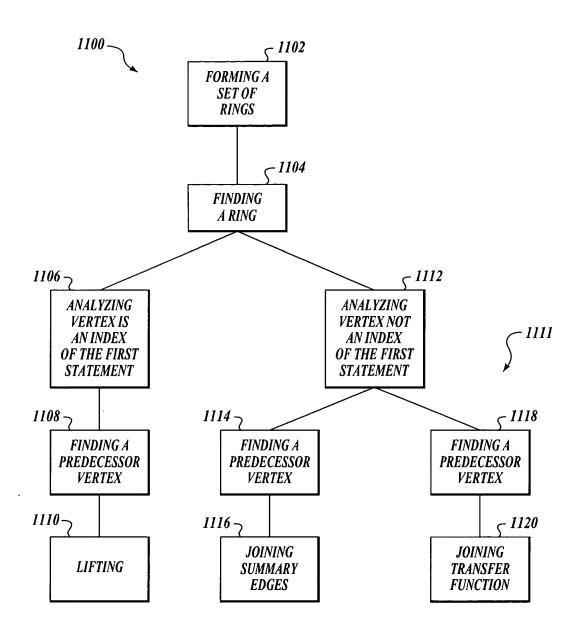


FIG. 11

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```
global
PE': V_G \rightarrow \text{set-of } (D \times D)
Worklist V_G \rightarrow \text{set-of}(D \times D)
procedure Propagate(V:V_{G,p}:(D \times D))
begin
  if p \notin PE'(v) then
      PE'(v) := PE'(v) \cup \{p\}
      Worklist(v) := Worklist(v) \cup \{p\}
  fi
end
procedure CMOP<sub>SP<sub>rhs</sub></sub>(S : set-of D)
begin
  PE'(entry) := \{\langle d, d \rangle \mid d \in S\}
   Worklist(entry) := PE'(entry)
   while \exists v_2 \ s.t \ Worklist(v_2) \neq 0 do
     select and remove \langle d_1, d_2 \rangle from Worklist(v_2)
     for each v_2 \rightarrow v_3 \in E_G do
        for each d_3 \in M(v_2 \rightarrow v_3)(\{d_2\}) do
           Propagate(v_3\langle d_1, d_3\rangle)
        od
     od
  od
end
```

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```
global
PE': V \rightarrow \text{set-of D} \times \text{set-of D}
Worklist: V_G \rightarrow \text{set-of D} \times \text{set-of D}
procedure Propagate(v : V_{G,p} : (\text{set-of D} \times \text{set-of D}))
begin
  if p \notin PE'(v) then
      PE'(v) := PE'(v) \cup \{p\}
      Worklist(v) := Worklist(v) \cup \{p\}
  fi
end
procedure CSMOP<sub>SPrhs</sub>(S': set-of (set-of D))
begin
  PE'(entry) := \{\langle S, S \rangle \mid S \in S'\}
      Worklist(entry) := PE' (entry)
      while \exists v_2 \ s.t \ Worklist(v_2) \neq 0 do
        select and remove \langle S_1, S_2 \rangle from Worklist(v_2)
        for each v_2 \rightarrow v_3 \in E_G do
           let S_3 = M(v_2 \rightarrow v_3)(S_2) in
              Propagate(v_3 \langle S_1, S_3 \rangle)
           ni
        od
     od
  end
```